

In this exam:

- The height of a singleton tree is defined as 0.
- The order of a node in a tree is defined as the number of its children.
- The order of a tree is defined as the maximum order among its nodes.
- Unless otherwise specified, all logarithms are taken with respect to base 2, i.e., \log_2 .

一、(20%)是非題 (每題 2 分，答錯每題倒扣 1 分至本大題 0 分止)

1. Given a connected graph $G = (V, E)$, if a vertex $v \in V$ is visited during level k of a breadth-first search from source vertex $s \in V$, then every path from s to v has length at most k .
2. Depth-first search will take $\Theta(V^2)$ time on a graph $G = (V, E)$ represented as an adjacency matrix.
3. In a weighted undirected graph $G = (V, E, w)$, breadth-first search from a vertex s finds single-source shortest paths from s (via parent pointers) in $O(V + E)$ time.
4. In a weighted undirected tree $G = (V, E, w)$, breadth-first search from a vertex s finds single-source shortest paths from s (via parent pointers) in $O(V + E)$ time.
5. In a weighted undirected tree $G = (V, E, w)$, depth-first search from a vertex s finds single-source shortest paths from s (via parent pointers) in $O(V + E)$ time.
6. Dijkstra's shortest-path algorithm may relax an edge more than once in a graph with a cycle.
7. Given a weighted directed graph $G = (V, E, w)$ and a source $s \in V$, if G has a negative-weight cycle somewhere, then the Bellman-Ford algorithm will always compute an incorrect result for some $\delta(s, v)$.
8. In a weighted directed graph $G = (V, E, w)$ containing no zero- or positive-weight cycles, Bellman-Ford can find a longest (maximum-weight) path from vertex s to vertex t .
9. In a weighted directed graph $G = (V, E, w)$ containing a negative-weight cycle, running the Bellman-Ford algorithm from s will find a shortest acyclic path from s to a given destination vertex t .
10. Given a weighted directed graph $G = (V, E, w)$ and a shortest path p from s to t , if we doubled the weight of every edge to produce $G' = (V, E, w')$, then p is also a shortest path in G' .

二、(64%)單選題 (每題 2 分，答錯不倒扣)

Please choose the best possible answer in each of the following questions. In questions that ask for time complexities, please choose the tightest bound.

11. In a binary heap of size n , the worst case time complexity of one insert operation is
(A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
12. In a binary heap of size n , the worst case time complexity of one find-min operation is
(A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
13. In a binary heap of size n , the worst case time complexity of one decrease-key operation is
(A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
14. In a binomial heap of size n , the worst case time complexity of $O(\log n)$ consecutive insert operations is
(A) $O(\log n)$ (B) $O(\log^2 n)$ (C) $O(n \log n)$ (D) $O(n \log^2 n)$ (E) $O(n^2 \log n)$
15. In a binomial heap of size n , the worst case time complexity of one delete-min operation is
(A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$

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16. The worst case time complexity of one union operation upon two binomial heaps of total size n is
 (A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
17. In a binomial heap of size n , the worst case time complexity of one decrease-key operation is
 (A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
18. In a Fibonacci heap of size n , the worst case time complexity of one find-min operations is
 (A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
19. In a Fibonacci heap of size n , the worst case time complexity of one delete-min operations is
 (A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
20. The worst case time complexity of one union operation upon two Fibonacci heaps of total size n is
 (A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
21. In a leftist heap of size n , the worst case time complexity of one delete-min operation is
 (A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
22. The worst case time complexity of one merge operation upon two leftist heaps of total size n is
 (A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
23. In a leftist heap of size n , the worst case time complexity of one decrease-key operation is
 (A) $O(1)$ (B) $O(\log^*n)$ (C) $O(\log n)$ (D) $O(n)$ (E) $O(n \log n)$
24. An AVL tree of 100 nodes has height at least
 (A) 5 (B) 6 (C) 7 (D) 8 (E) 9
25. A 2-3-4 tree of 255 nodes has height at least
 (A) 3 (B) 4 (C) 5 (D) 6 (E) 7
26. How many binomial trees are there in a binomial heap with 2020 items?
 (A) 1 (B) 3 (C) 5 (D) 7 (E) 9
27. How many items are there in a binomial tree with order 4?
 (A) 13 (B) 14 (C) 15 (D) 16 (E) 17
28. Which of the following is correct?
 (A) $\frac{n^2}{\log n} = O(n^{1.5})$
 (B) $2^{2n} = \Theta(2^n)$
 (C) If $f(n) = O(g(n))$ and $g(n) = O(h(n))$, then $f(n) = O(h(n))$
 (D) If $f(n) \neq O(g(n))$, then $g(n) = O(f(n))$
29. Order the following functions by the big-Oh notation in decreasing order.
 (a) 1.001^n (b) $n!$ (c) $n^{5/2}$ (d) $2^{7 \log n}$ (e) 3^{100000} (f) $\log \log n$ (g) \sqrt{n} (h) n^n
 Which of the following order is correct?
 (A) $h > a > b > d > c > g > f > e$
 (B) $h > b > a > d > c > g > f > e$
 (C) $h > b > a > d > c > f > g > e$
 (D) $b > h > a > d > c > g > f > e$
30. If $T(n) = 2T(\sqrt{n}) + \log n$, where $T(2) = 0$, then $T(n) = ?$
 (A) $n/\log n$
 (B) $\log n (\log \log n)$
 (C) $\log \log n$
 (D) $\log n$
31. Which is the tight lower bound for any comparison based sorting algorithms?
 (A) $\Omega(n)$
 (B) $\Omega(n^2)$
 (C) $\Omega(n \log n \log n)$
 (D) $\Omega(n \log n)$

32. What's the time complexity to do (deterministic) quick sort, insertion sort and merge sort on a list of numbers in perfect sorted order?

- (A) quick sort: $O(n^2)$, insertion sort: $O(n)$, merge sort: $O(n \log n)$
 (B) quick sort: $O(n \log n)$, insertion sort: $O(n^2)$, merge sort: $O(n \log n)$
 (C) quick sort: $O(n)$, insertion sort: $O(n^2)$, merge sort: $O(n \log n)$
 (D) quick sort: $O(n^2)$, insertion sort: $O(n^2)$, merge sort: $O(n)$

33. To search a specific number in a list, there are following two strategies:

- (a) Linear Search
 (b) First sort the list using merge sort, and then do binary search

Which strategy is faster in average? What's its corresponding time complexity?

- (A) (a), $O(n)$
 (B) (b), $O(\log n)$
 (C) (a), $O(\log n)$
 (D) (b), $O(n \log n)$

34. Consider the following list of 9 activities with their start time s and end time t , what's the maximum number of activities that can be performed by a single person? Assuming that a person can only work on a single activity at a time.

index	1	2	3	4	5	6	7	8	9
s	2	3	1	10	14	6	13	7	19
t	5	9	12	16	17	15	18	11	20

- (A) 2
 (B) 3
 (C) 4
 (D) 5

35. Consider the following two sequences:

3, 2, 1, 7, 9, 8, 6, 8, 0, 9

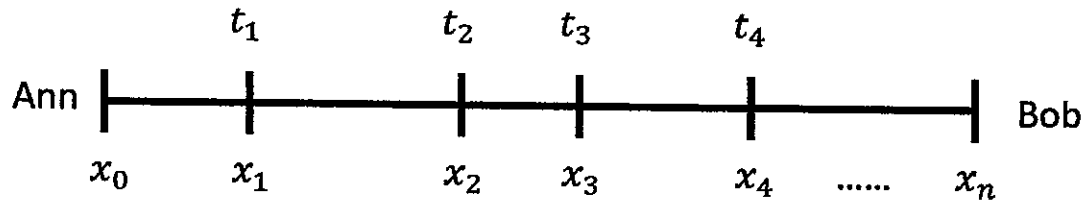
8, 3, 7, 8, 1, 9, 8, 3, 9, 0

The length of the longest common subsequence is:

- (A) 3
 (B) 5
 (C) 4
 (D) 2

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For problems 36 to 42, please refer to the following dynamic programming problem.



Ann is riding scooter from her home to Bob's home with a constant speed v . However, her scooter can only hold gas for 100kms riding at most. There are a few gas stations along the way at x_1, x_2, \dots, x_n distances from her home, but filling up at different stations x_i would take different time t_i . She would start her trip with a full tank of gas, and if she decides to fill up at a gas station, she needs to fill her entire tank up. Since the riding time is constant (x_n/v), we only need to consider the time consumed by filling gas. (Assume that the distances between two consecutive stations are not longer than 100kms. Bob's home is at x_n distances from Ann's home and the distances between his home and the last station are not longer than 100kms as well.)

36. If $s[a]$ indicates the minimum time she has spent on filling gas when she has filled the gas at a^{th} station, then

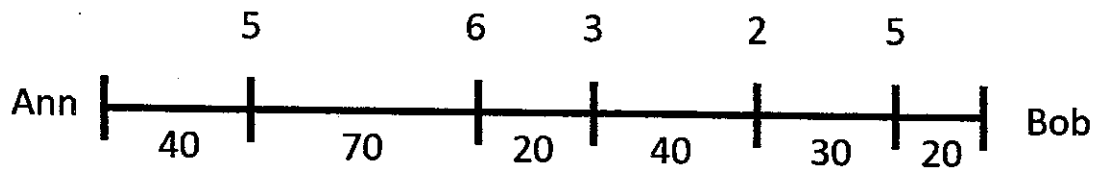
$$(A) \quad s[a] = \begin{cases} 0, & \text{if } a = 0 \\ \min_{(0 \leq i \leq a-1) \wedge (x_a - x_i \leq 100)} (s[i] + t[i]), & \text{otherwise} \end{cases}$$

$$(B) \quad s[a] = \begin{cases} 0, & \text{if } a = 0 \\ \min_{(0 \leq i \leq a-1) \wedge (x_a - x_i \leq 100)} (s[i] + t[a]), & \text{otherwise} \end{cases}$$

37. If $r[a]$ indicates the minimum time she has spent on filling gas when she arrives at a^{th} gas station, then

$$(A) \quad r[a] = \begin{cases} 0, & \text{if } a = 0 \\ \min_{(0 \leq i \leq a-1) \wedge (x_a - x_i \leq 100)} (r[i] + t[i]), & \text{otherwise} \end{cases}$$

$$(B) \quad r[a] = \begin{cases} 0, & \text{if } a = 0 \\ \min_{(0 \leq i \leq a-1) \wedge (x_a - x_i \leq 100)} (r[i] + t[a]), & \text{otherwise} \end{cases}$$



According to the figure, solve $s[1], \dots, s[5]$.

38. $s[1] =$

- (A) 0
- (B) 5
- (C) 6
- (D) 40

39. $s[2] =$

- (A) 0
- (B) 5
- (C) 6
- (D) 11

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40. $s[3] =$

- (A) 5
- (B) 8
- (C) 11
- (D) 13

41. $s[4] =$

- (A) 8
- (B) 10
- (C) 11
- (D) 13

42. $s[5] =$

- (A) 8
- (B) 10
- (C) 11
- (D) 13

三、(16%) 複選題 (每題 2 分，每答錯一個選項倒扣 0.4 分至該題 0 分止)

43. Consider the expression tree in Fig.1

- (A) The prefix expression is $a+b*c+d+e*f*g*h+i+j*k$
- (B) The prefix expression is $*++a*bc*+def*g++hi*jk$
- (C) The infix expression is $*+++g+a*f+*bcdehijk$
- (D) The postfix expression is $abc*+de+f*+ghi+jk*+**$
- (E) The postfix expression is $k*j+i+h*g*f*e+d+c*b+a$

44. In a splay tree of size n

- (A) The height is $O(\log n)$.
- (B) Time complexity of one insertion is $O(\log n)$.
- (C) Time complexity of n insertions is $O(n \log n)$.
- (D) Time complexity of one deletion is $O(\log n)$.
- (E) Time complexity of n deletions is $O(n \log n)$.

45. Consider the splay tree in Fig.2.

- (A) In Fig.2, after inserting the key 93, then the parent of key 83 is key 79.
- (B) In Fig.2, after inserting the key 80, then the parent of key 80 is key 79.
- (C) In Fig.2, after inserting the key 33, then the parent of key 35 is key 48.
- (D) In Fig.2, after inserting the key 1, then the parent of key 15 is key 2.
- (E) In Fig.2, after inserting the key 40, then the root is key 40.

46. In a red-black tree of size n

- (A) Time complexity of one search is $O(\log n)$.
- (B) Time complexity of one insertion is $O(\log n)$.
- (C) Time complexity of n insertions is $O(n \log n)$.
- (D) Time complexity of one deletion is $O(\log n)$.
- (E) Time complexity of n deletions is $O(n \log n)$.

47. Which of the followings are balanced trees?

- (A) Binary search tree.
- (B) AVL tree.
- (C) Red-black tree.
- (D) Splay tree.
- (E) 2-3-4 tree.

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48. In an AA tree, which of the following statements are correct?

- (A) There are no two consecutive right horizontal links.
- (B) There is no left horizontal link.
- (C) A node has the same level as its right horizontal link child.
- (D) A node has the same level as its left-link child.
- (E) The height is $O(\log n)$.

49. Which of the following operations are supported by binary search trees?

- (A) Find key
- (B) Insert key
- (C) Delete key
- (D) Find-Min
- (E) Delete-Min

50. Consider the Fibonacci heap in Fig.3

- (A) In Fig.3, after decreasing key 89 to 40, there will be 3 marked nodes in total.
- (B) In Fig.3, after decreasing key 59 to 40, there will be 5 marked nodes in total.
- (C) In Fig.3, after decreasing key 15 to 10, there will be 2 trees in total.
- (D) In Fig.3, after delete-min, there will be two trees.
- (E) In Fig.3, after inserting key 68, there will be one tree.

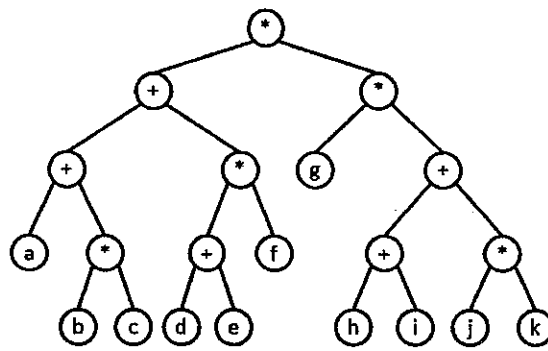


Fig.1 Expression tree

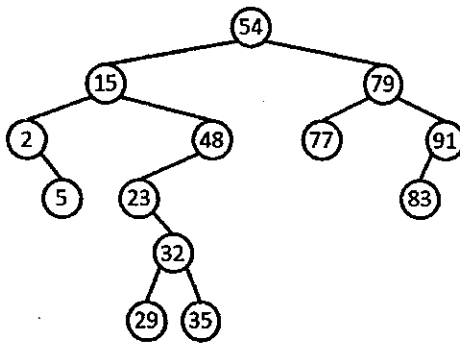


Fig.2 Splay tree

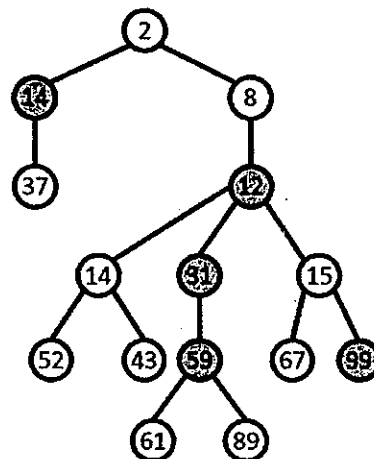


Fig.3 Fibonacci heap